

# Competing Inheritance Paths in Dependent Type Theory

a case study in functional analysis

Reynald Affeldt, Cyril Cohen, Marie Kerjean,  
Assia Mahboubi, Damien Rouhling, Kazuhiko Sakaguchi

AIST / Inria / U. Tsukuba

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# What this Presentation is About

## inheritance in hierarchies of structures

- Hierarchies of structures are important to organize formalizations
- They are implemented with type classes or unification hints
- Problems for developers: performance/predictability,  
not all issues are properly documented
- In this talk, we discuss problems caused by competing definitions
- In particular, we identify **forgetful inheritance** as a solution to  
the problem of competing inheritance paths for poorer structures

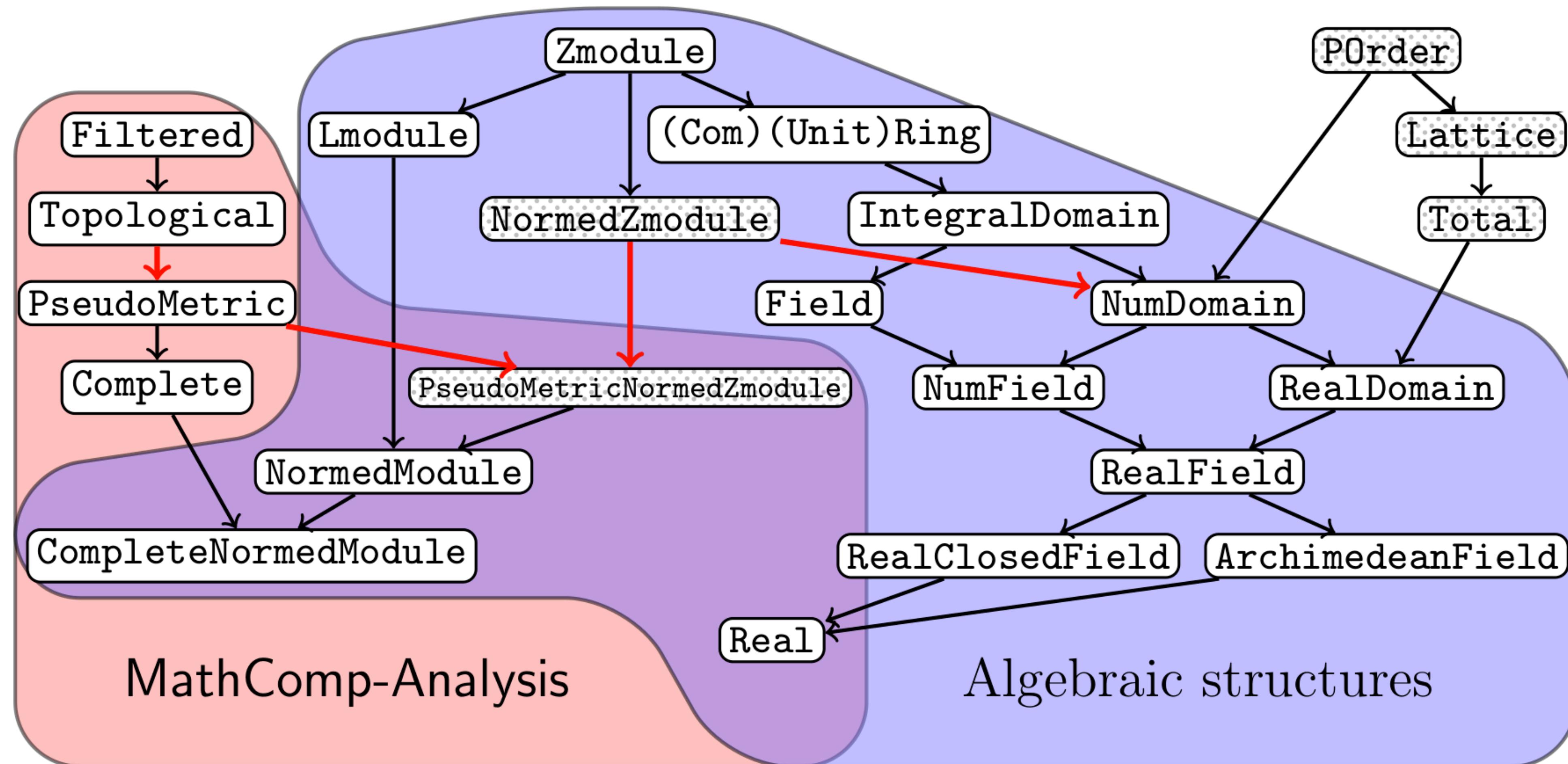
# Motivation for this Work

## analysis with the Coq proof assistant

- **MathComp** is the library for algebra used to formalize the odd order theorem by Gonthier et al. [ITP 2013]
- **Coquelicot** is a library for real analysis by Boldo et al. [MCS 2015]
  - It extends and improves the standard library of Coq
- **MathComp-Analysis** is a work-in-progress that extends MathComp for functional analysis.
  - It is inspired by Coquelicot but comes with original features [JFR 2018]

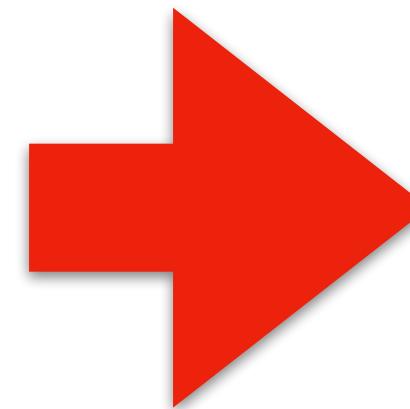
# MathComp-Analysis Hierarchy

the concrete result of this presentation (teaser)



# Competing Inheritance Paths in Dependent Type Theory

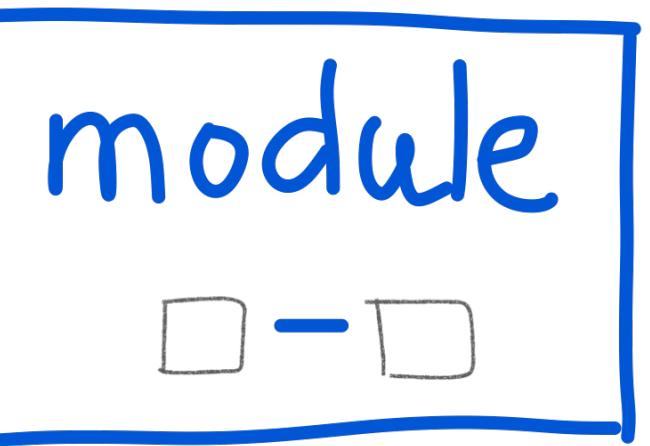
## outline



1. Background: Hierarchies of Structures
2. Problem: Extend MathComp with Analysis
3. Forgetful Inheritance for Competing Inheritance Paths
4. Application to MathComp-Analysis
5. Conclusion

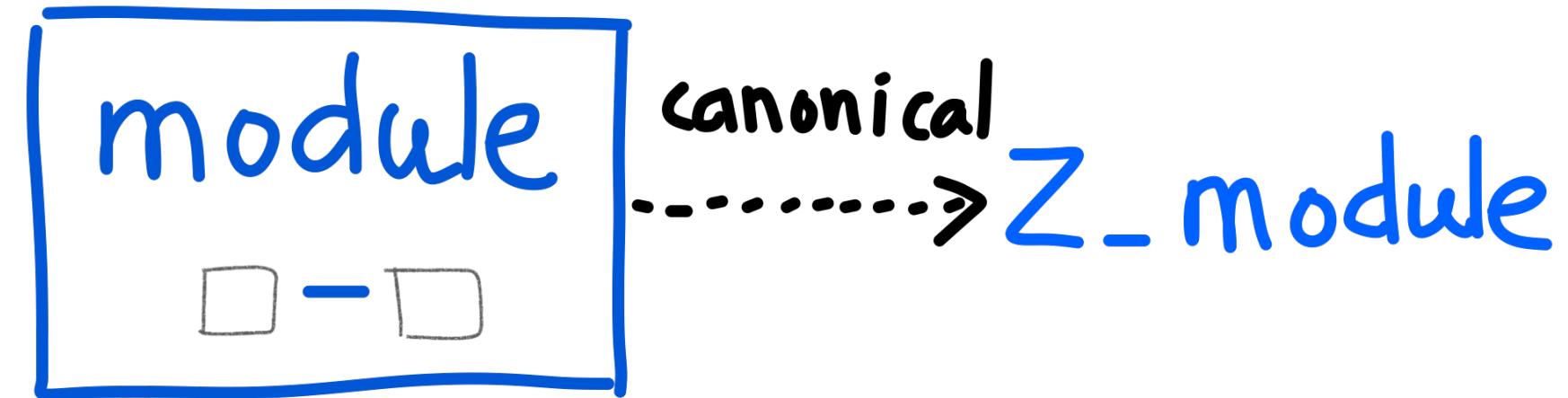
# Mathematical Structure

using a packed class [Garillot et al., 2009]



- 1 Record **isModule**  $T := \text{IsModule} \{ \text{minus\_op} : T \rightarrow T \rightarrow T \}$ . } **Mixin** = operator + properties
- 2 Structure **module** := Module {  
  **module\_carrier** : Type ;  
  **module\_isModule** : isModule module\_carrier }. } **Structure** = carrier + mixin
- 3 Definition **minus** ( $M : \text{module}$ )  
  : module\_carrier M  $\rightarrow$  module\_carrier M  $\rightarrow$  module\_carrier M } **operator lifted**  
  := minus\_op \_ (module\_isModule M). } from the **mixin** to the **structure**
- 4 Notation “ $x - y$ ” := (minus \_ x y). } **Notation with a hole**  $\approx$  overloading

# Structure Inference with unification hints



- 1 Definition  $Z\_module := \text{Module } Z (\text{IsModule } Z Z.\text{sub})$ . } structure instance
- 2 Canonical  $Z\_module$ . } the instance is canonical

Check forall  $x y : Z, x - y = x - y$ . Unset  
Printing  
Notations  $\rightarrow$  minus  $\underline{?} \ x \ y$

$x$  and  $y$  are both annotated with  $Z$ .  
 $x$  is annotated with "module".  
 $y$  is annotated with "module\_carrier ?".  
A blue arrow points from  $x$  to  $y$ , labeled "type equation".

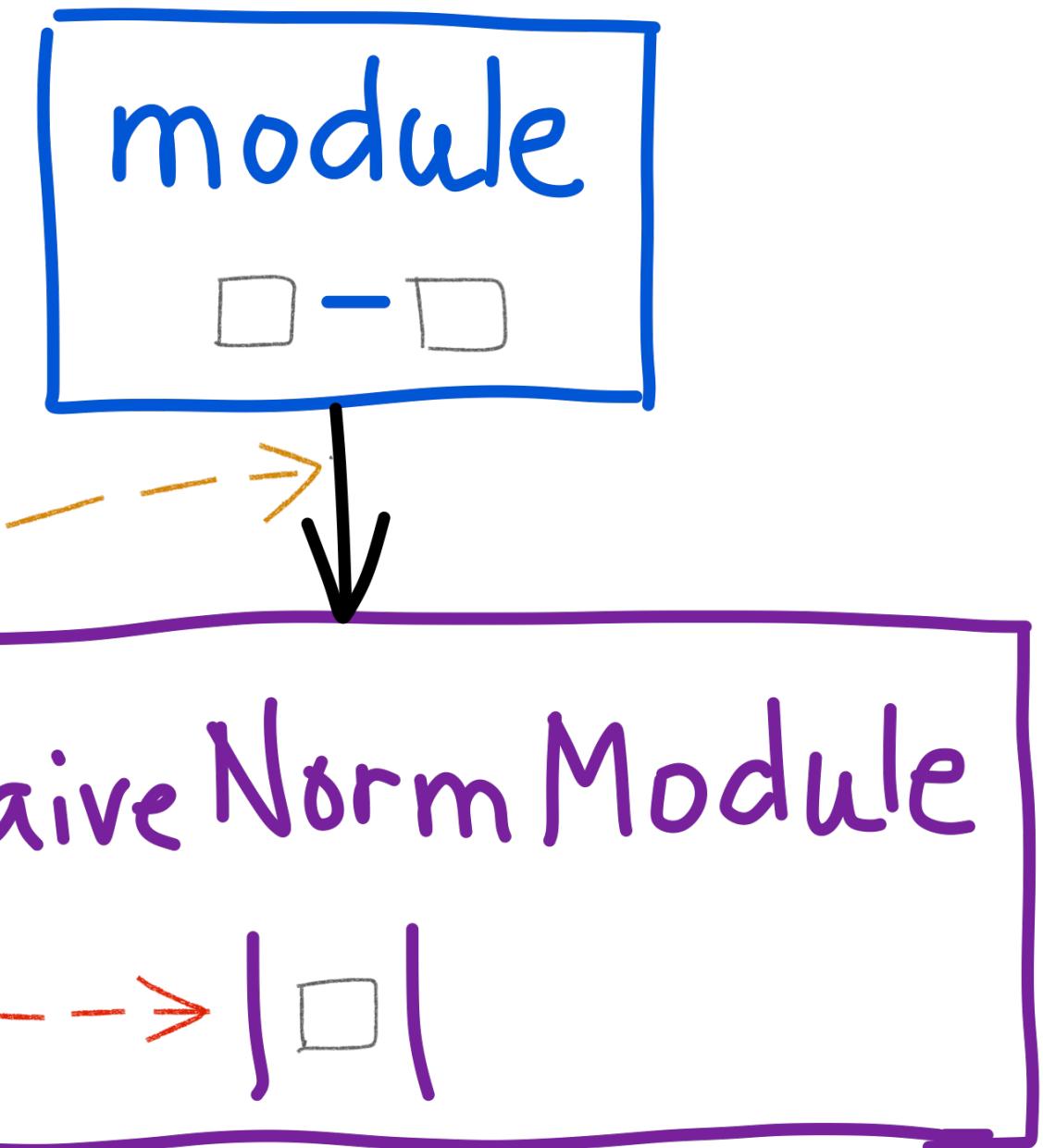
$? \equiv Z\_module$  solution  $\leftarrow \dots$  module\_carrier ?  $\equiv Z$

# Inheritance of Structures using the notion of class

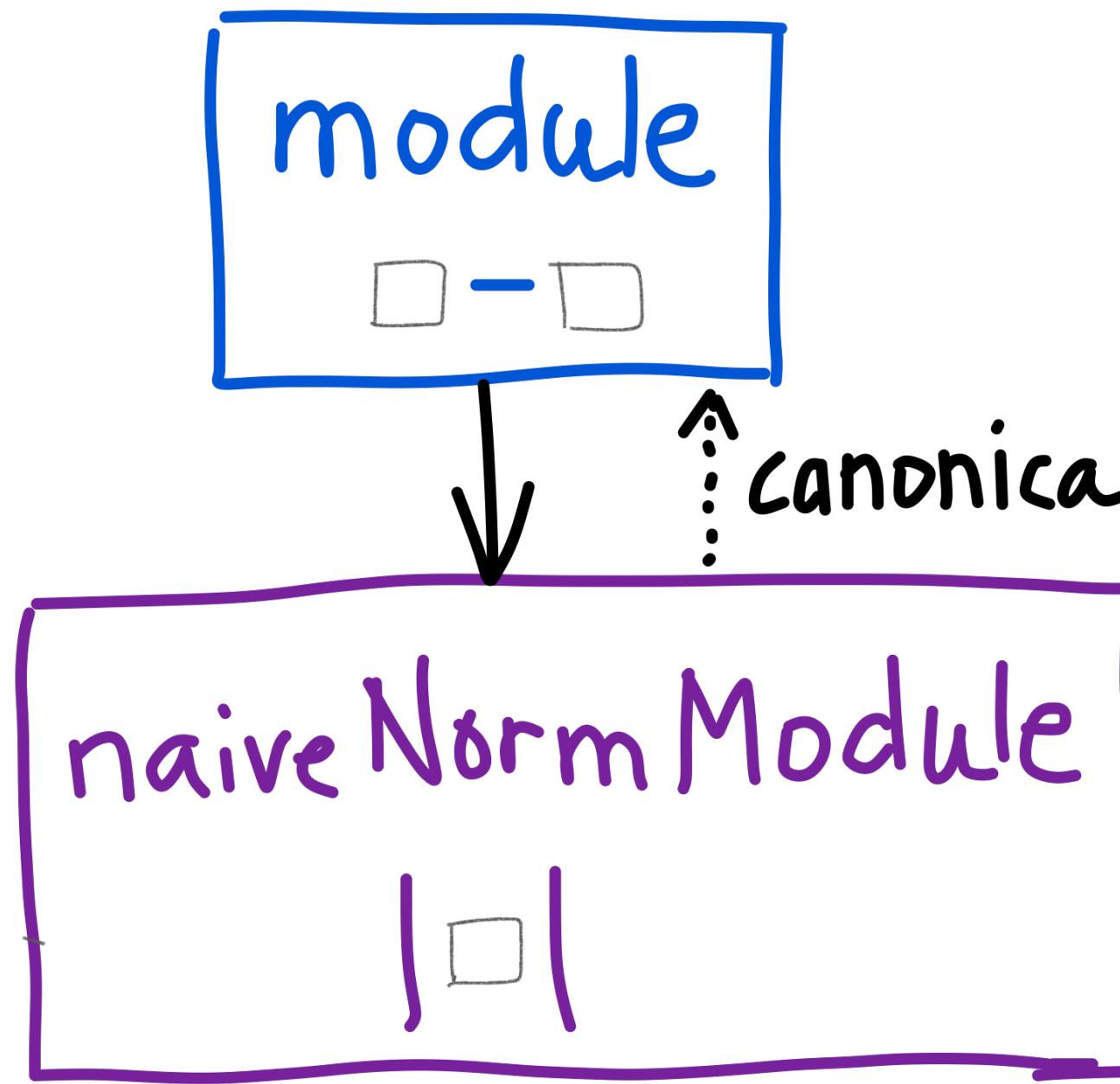
1 Record **naiveNormMixin** (*T* : module) := NaiveNormMixin { } **mixin**  
*naive\_norm\_op* : *T* -> nat }.

Record **isNaiveNormModule** (*T* : Type) := IsNaiveNormModule { }  
*nbase* : isModule *T* ;  
*nmix* : naiveNormMixin (Module \_ *nbase*) }.

2 Structure **naiveNormModule** := NaiveNormModule { } **structure**  
*naive\_norm\_carrier* :> Type ;  
    implicit coercion  
*naive\_normModule\_isNormModule* : isNaiveNormModule *naive\_norm\_carrier* }.



# Inference in Presence of Inheritance using unification hints



Check forall (N : naiveNormModule) (x y : N),  
module\_carrier ?  $\xrightarrow{x - y = x - y}$   $\text{naive\_norm\_carrier } N$   
KO *← implicit coercion*

Canonical naiveNorm\_isModule (N : naiveNormModule) :=  
Module N (nbase \_ (naive\_normModule\_isNormModule N)).

module\_carrier ?  $\equiv$  naive\_norm\_carrier N  
OK  
naiveNorm\_isModule N

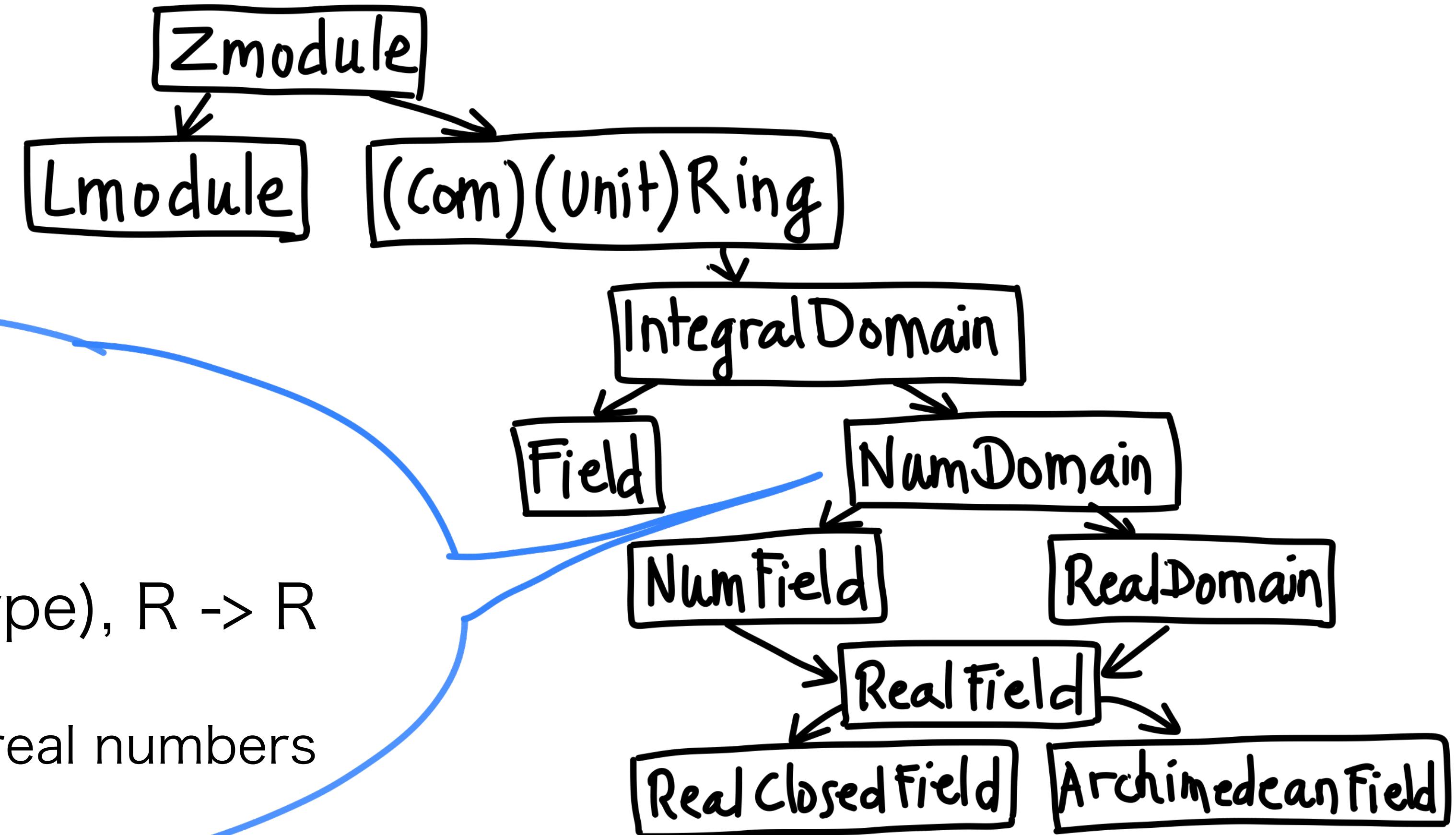
# Competing Inheritance Paths in Dependent Type Theory

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# The MathComp Library

(excerpt, before this work)



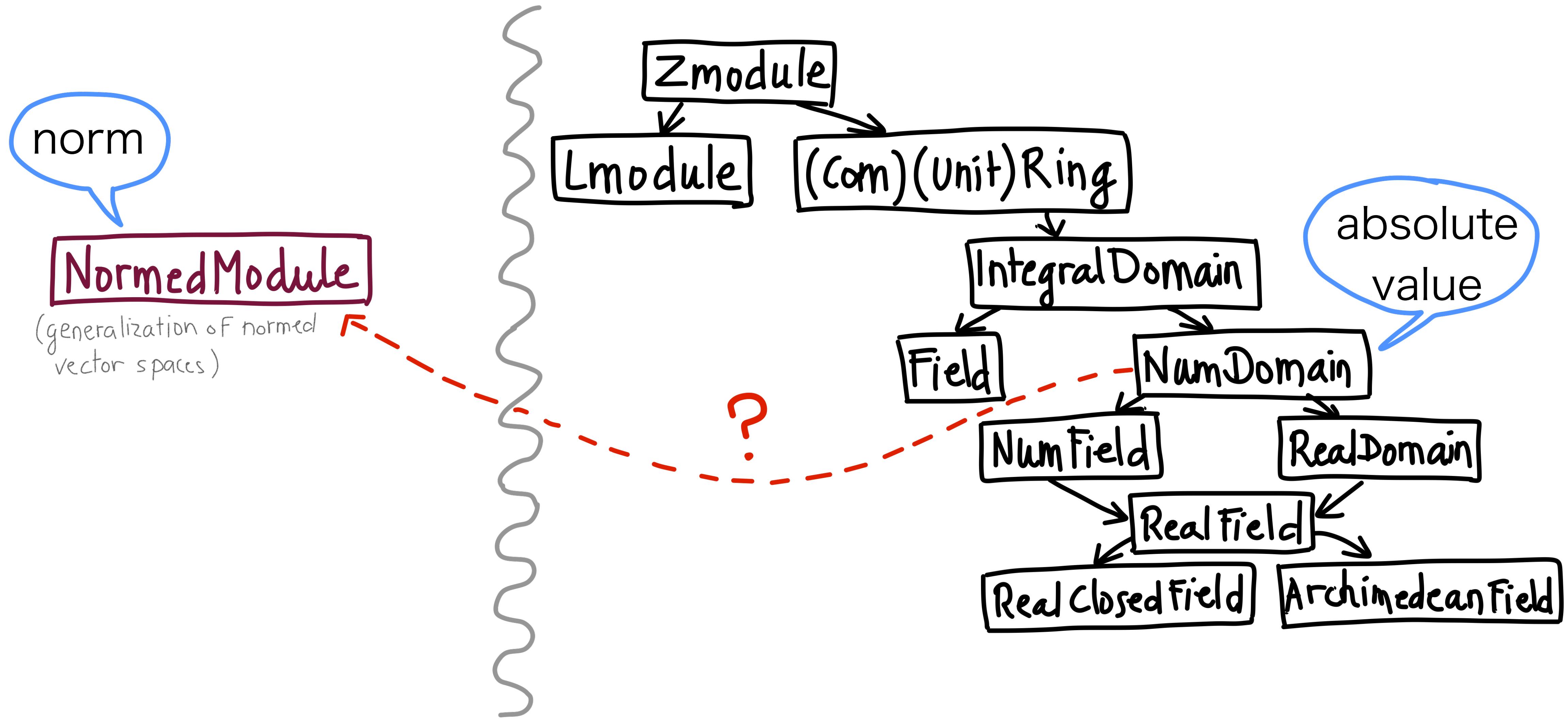
absolute value of type:

forall (R : numDomainType), R -> R

Sample instances: integers, real numbers

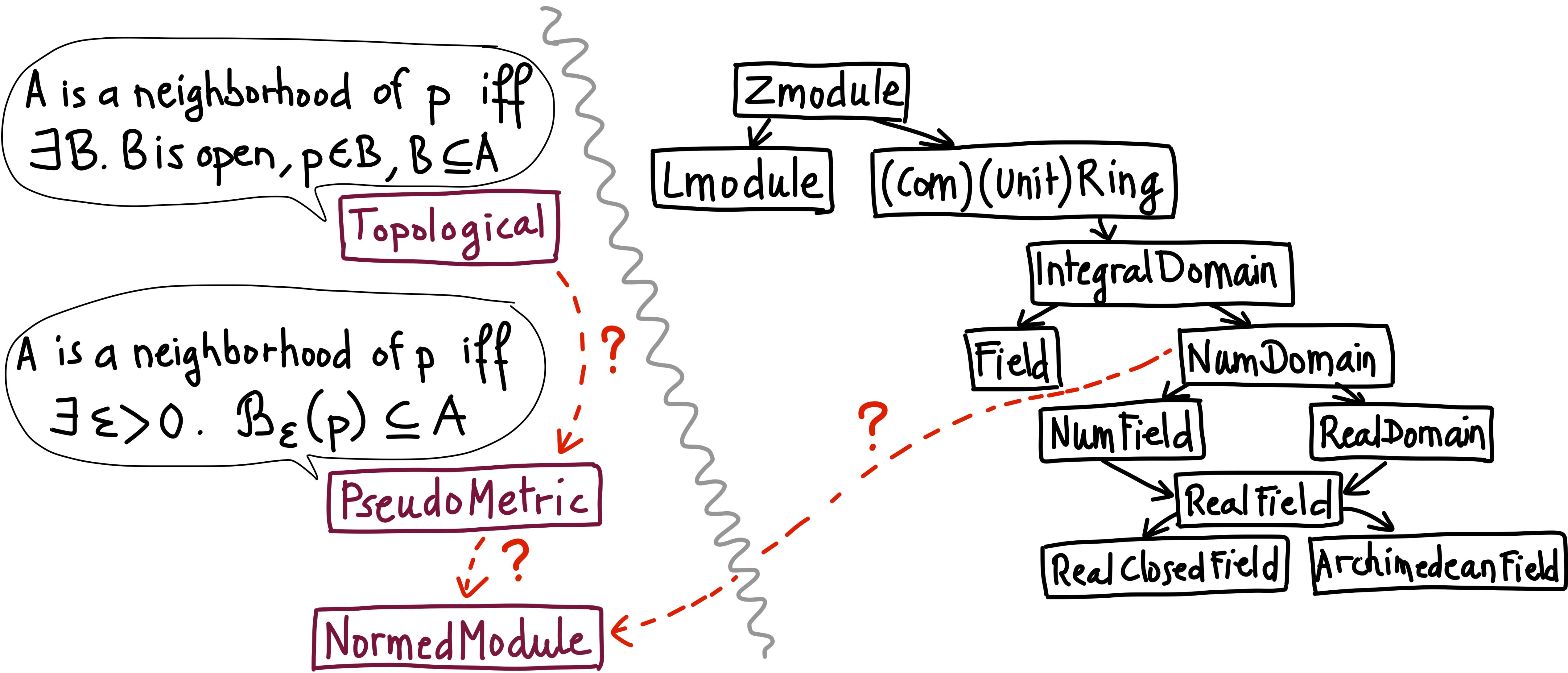
# Towards Analysis with MathComp

## issue #1: properties of differentiability, continuity



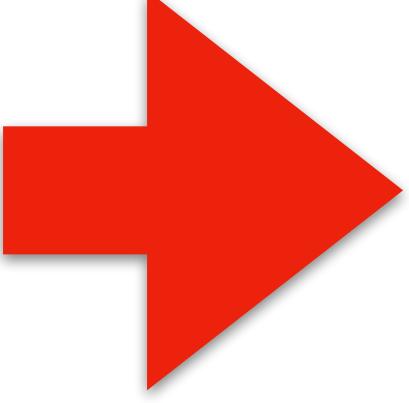
# Towards Analysis with MathComp

## issue #2: topology



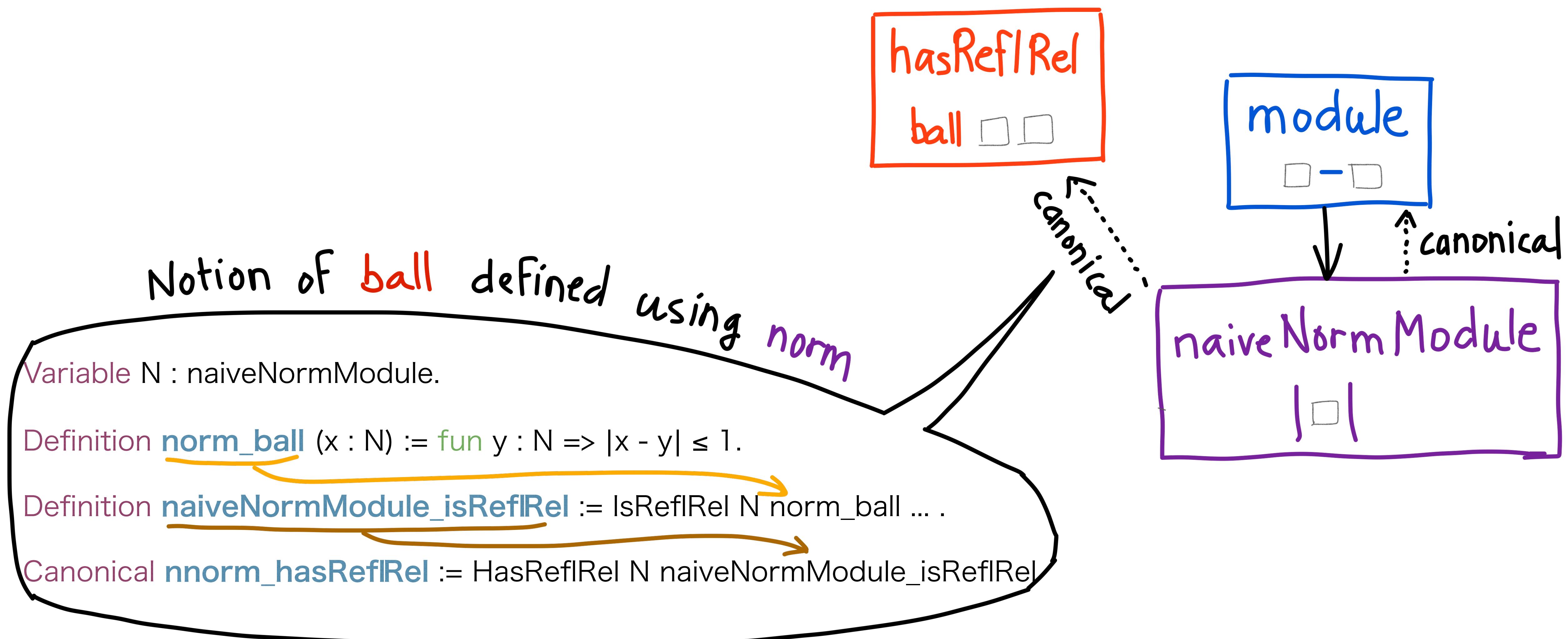
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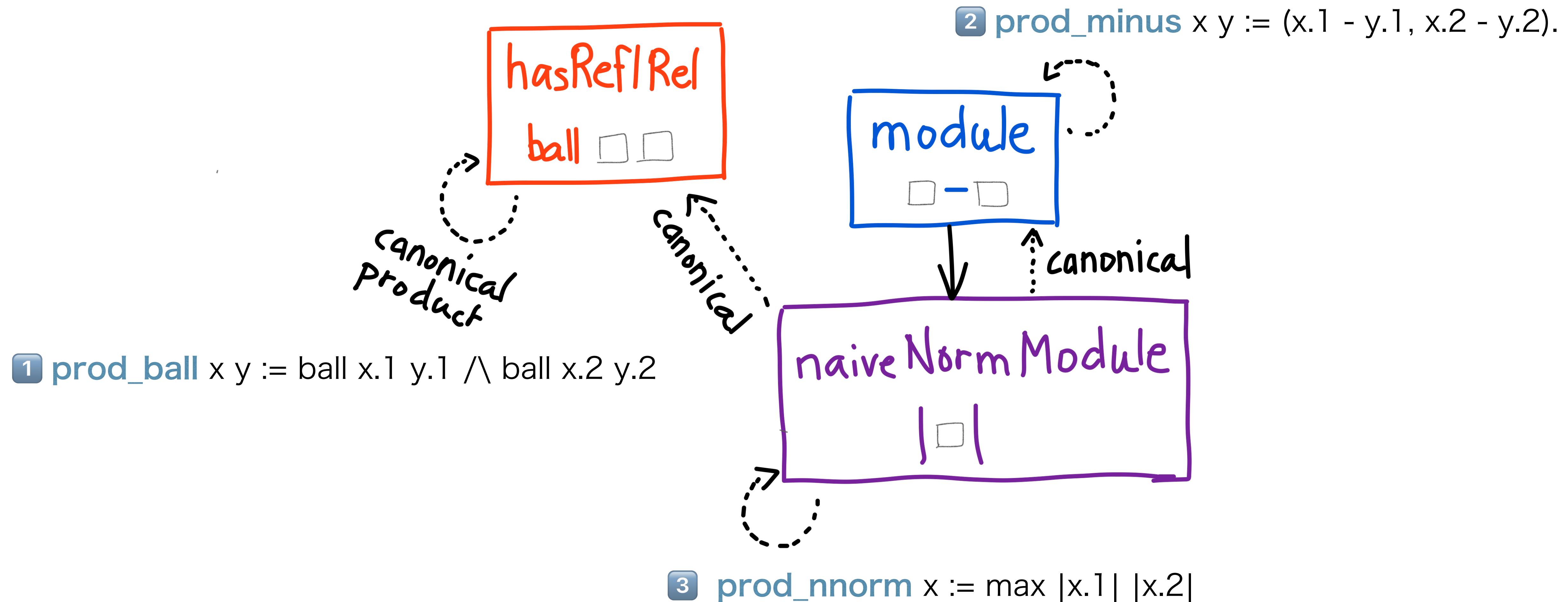
# Hierarchy Extension

a new structure with a notion of ball



# Hierarchy Extension

## canonical products for all structures



# A Puzzling Inference Failure

Example failure

*flyothesis* → (Pball : forall V : naiveNormModule, forall v : V, P (ball v))

(W : naiveNormModule) (w : W \* W)

*Goal* → : P (ball w).

Proof. Fail apply Pball. Abort.



Shouldn't this succeed since  
there is a canonical way to  
build products? 🤔

# A Puzzling Inference Failure

Example failure

nnorm\_hasRefIRel ?

prod\_naiveNormModule WW

(Pball : forall V : naiveNormModule, forall v : V, P(ball v))

(W : naiveNormModule) (w : W \* W)

: P(ball w).

Proof. Fail apply Pball. Abort.

prod\_hasRefIRel (nnorm\_hasRefIRel W)  
(nnorm\_hasRefIRel W)

# A Puzzling Inference Failure

Example failure

(Pball : forall V : naiveNormModule, forall v : V, P(ball v))

(W : naiveNormModule) (w : W \* W)

: P(ball w).

Proof. Fail apply Pball. Abort.

prod\_naiveNormModule WW

norm\_hasReflRel ?

ball\_op xy  $\triangleq |x-y| \leq 1$

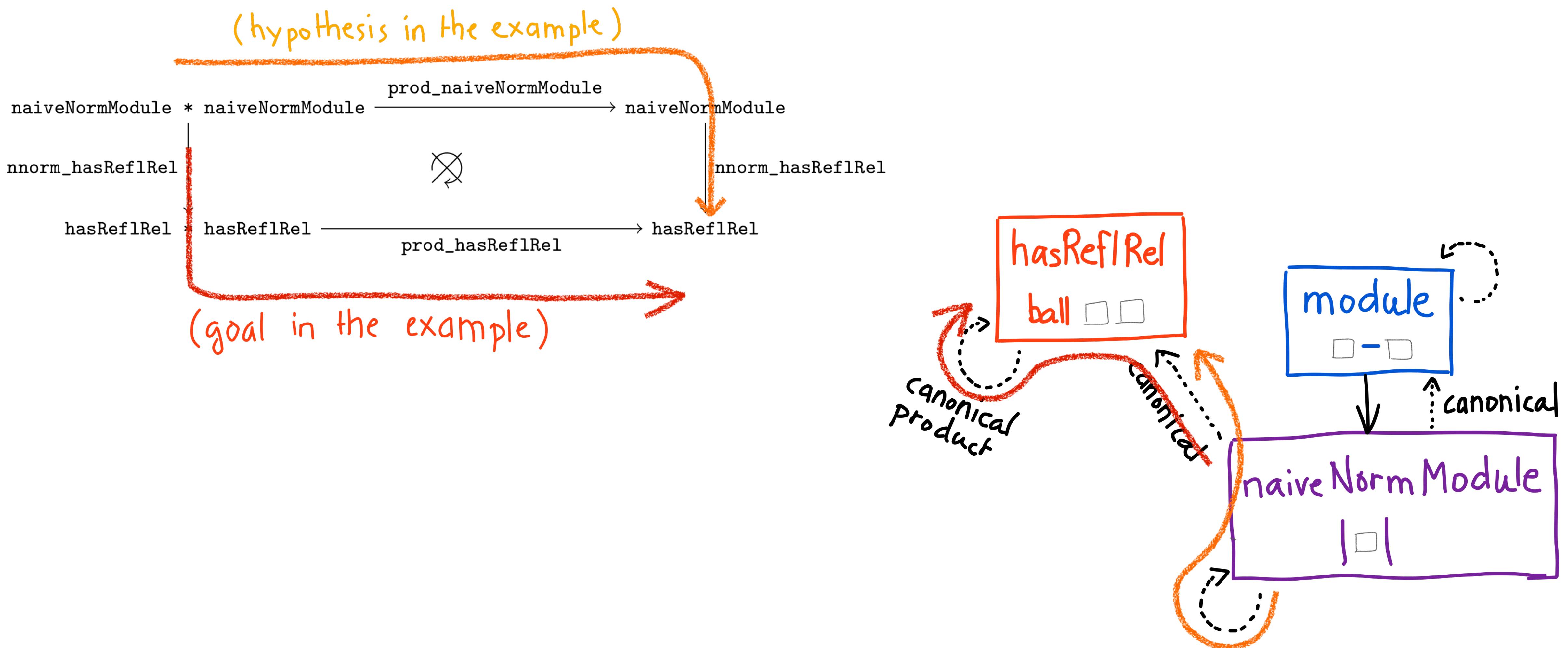
||| KO

not  
definitionally  
equal

ball\_op xy  $\triangleq \text{ball } x_1, y_1 \wedge \text{ball } x_2, y_2$

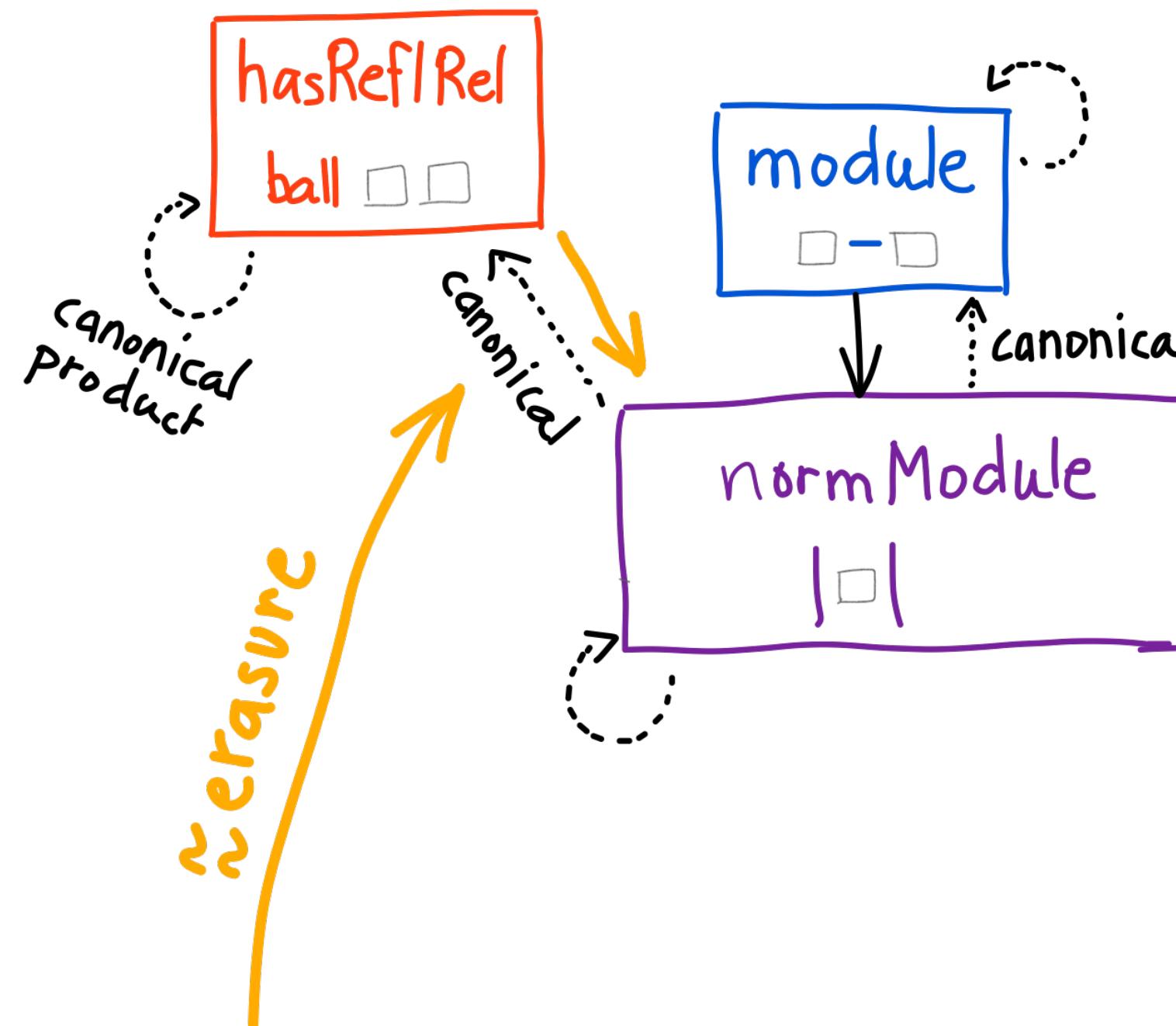
prod\_hasReflRel (norm\_hasReflRel W)  
(norm\_hasReflRel W)

# Problem of Competing Inheritance Paths diagrammatically



# Forgetful Inheritance

## how to fix the hierarchy



4 Canonical **norm\_hasRefIRel** (N : normModule) :=

HasRefIRel N (bmix \_ (normModule\_isNormModule N))

- hasRefI Rel as a parameter**
- 1 Record **normMixin** (T : module) (m : isRefIRel) := NormMixin {
- norm\_op : T -> nat ;
- norm\_ball\_opP : forall x y, ball\_op \_ m x y <-> norm\_op (x - y) ≤ 1 }.
- 2 Record **isNormModule** (T : Type) := IsNormModule {
- base : isModule T ;
- bmix : isRefIRel T ;
- mix : normMixin (Module \_ base) bmix }.
- 3 Structure **normModule** := NormModule {
- norm\_carrier :> Type ;
- normModule\_isNormModule : isNormModule norm\_carrier }.

# Competing Inheritance Paths with Type Classes

same problem, same solution

Type classes in Lean

Type classes in Coq

Packed classes in Coq



<https://math-comp.github.io/competing-inheritance-paths-in-dependent-type-theory/>

```
1 From Coq Require Import ssreflect ZArith.
2 Reserved Notation "| x |" (at level 30).
3 Reserved Notation "x ~~ y" (at level 30).
4
5 (* **** *)
6 (* Section 2.1 *)
7 (* **** *)
8 Record flatNormSpace := FlatNormSpace {
9   carrier : Type ;
10  fminus : carrier -> carrier -> carrier;
11  fnorm : carrier -> nat;
12  fnormP : forall x : carrier, fnorm (fminus x x) = 0}.
13
14 Check fminus.
15 (* fminus : forall f : flatNormSpace,
16   (* carrier f -> carrier f -> carrier f *) *)
17 Check fnormP.
18 (* fnormP : forall (f : flatNormSpace) (x : carrier f), *)
19 (*   fnorm f (fminus f x x) = 0 *)
20
21 Lemma Z_normP (n : Z) : Z.abs_nat (Z.sub n n) = 0.
22 Proof. by rewrite Z.sub_diag. Qed.
23 Definition Z_flatNormSpace := FlatNormSpace Z Z.sub Z.abs_nat Z_normP.
24
25 (* **** *)
26 (* Section 2.2 *)
27 (* **** *)
28 Record isModule T := IsModule {minus_op : T -> T -> T}.
29 Structure module := Module {
30   module_carrier : Type ;
31   module_isModule : isModule module_carrier}.
32
33 Definition minus {M : module} := minus_op _ (module_isModule M).
34 Notation "x - y" := (minus x y).
35
36 Definition Z_isModule : isModule Z := IsModule Z Z.sub.
37 Definition Z_module := Module Z Z_isModule.
38
39 Fail Check forall x y : Z, x - y = x - y.
40
41 Canonical Z_module.
42 Check forall x y : Z, x - y = x - y.
43 Print Canonical Projections.
44 (* the line Z <- module_carrier (Z_module) tries to solve *)
45 (* the equation `module_carrier ?M =?= Z` *)
46
47 (* begin omitted from the paper *)
48 Coercion module_carrier : module ->- Sortclass.
49 (* end omitted from the paper *)
```

Section 2.2  
isModule is in fact has\_sub --  
-- Section 2.3 --  
universe u

v :=

inu : T -> T -> T].

es.

inus : T -> T -> T].

Module Z.sub.

t.

Module(isModule Z, id 0)

giveNormSpace {

ve\_norm (x - x) = 0}.

\_instances.

y : N), x - y = x - y.

n n) = 0.

nSpace Z :=

(v \* v') := (

ace {

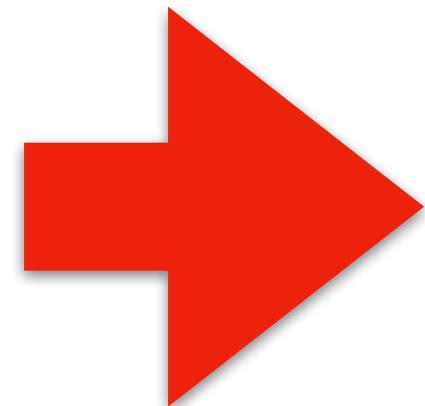
ances.

alpha, P (ball v))

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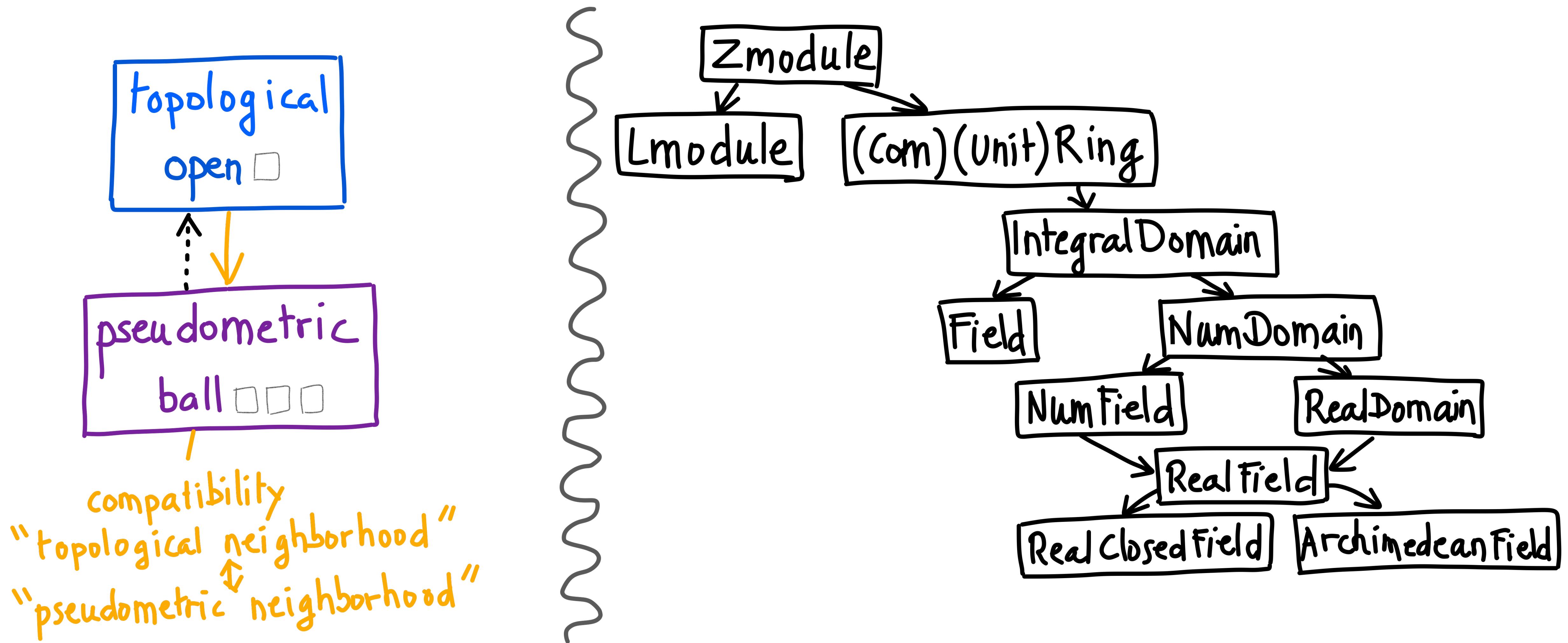
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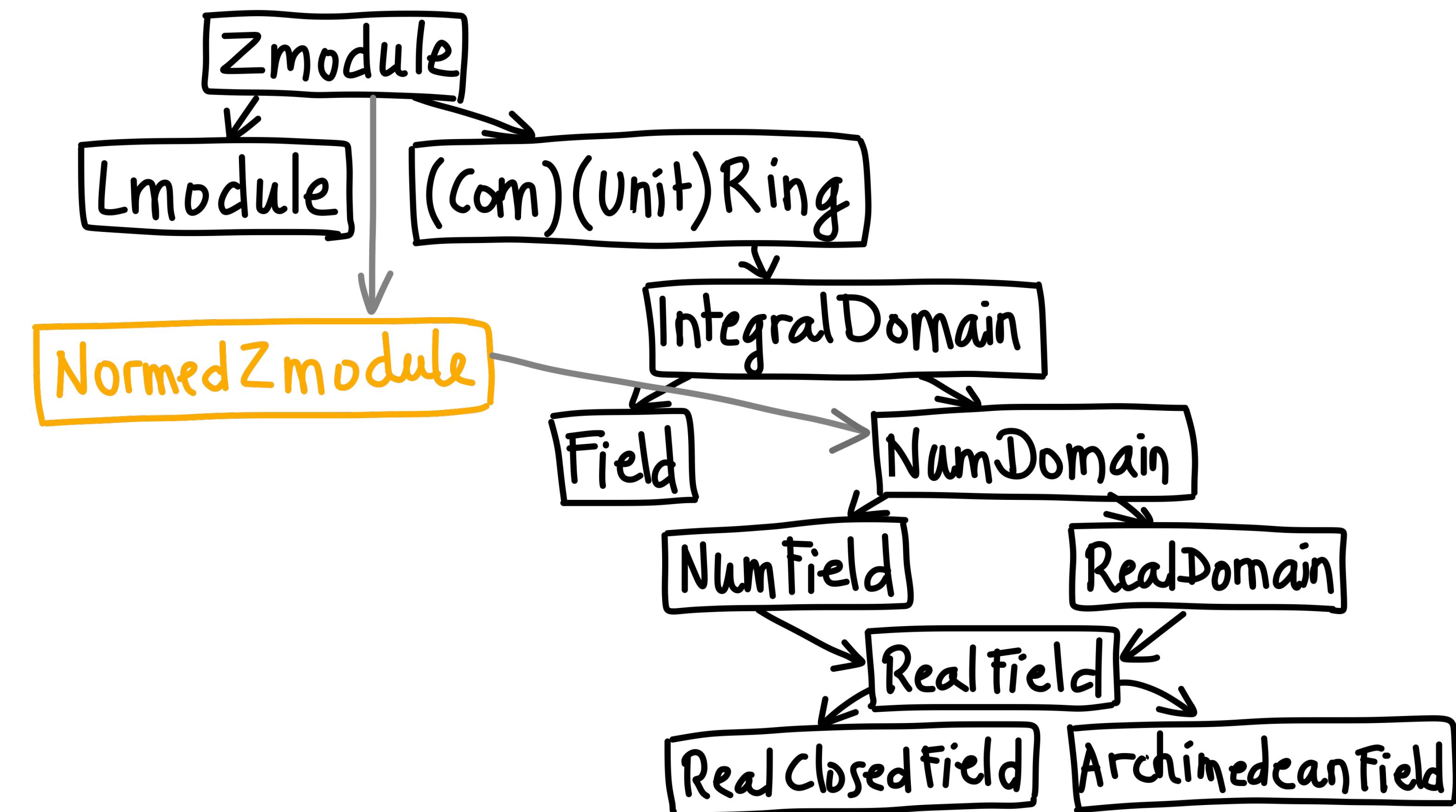
# Forgetful Inheritance

from pseudometric spaces to topological spaces



# Unify Absolute Value and Norm

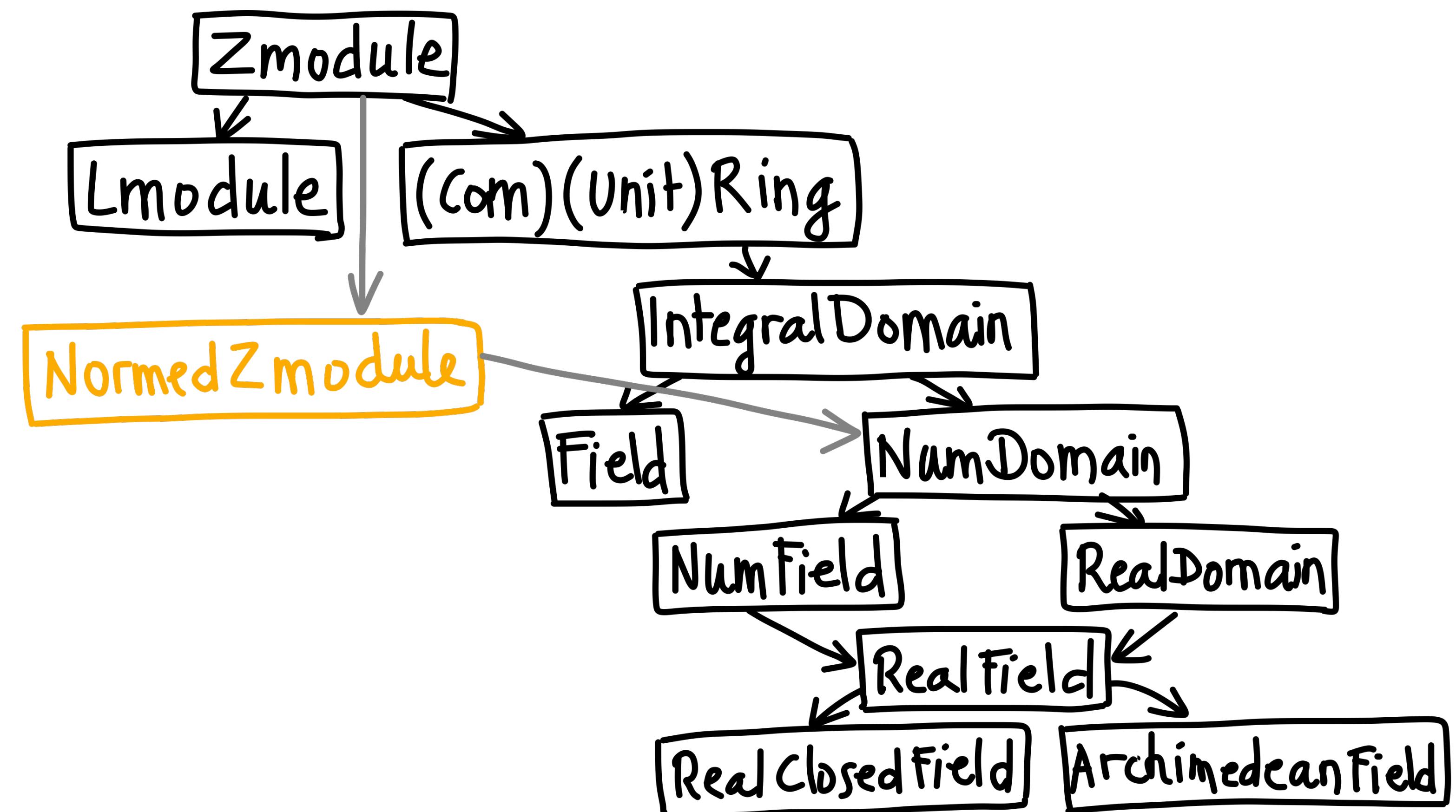
step #1: factorize with normed Abelian groups



# Unify Absolute Value and Norm

step #2: solve mutual dependency problem

properties of the norm require the codomain of the norm to be a normed Abelian group and an order (e.g., triangle inequality)...



# Forgetful Inheritance

## from numerical domain to normed Abelian group

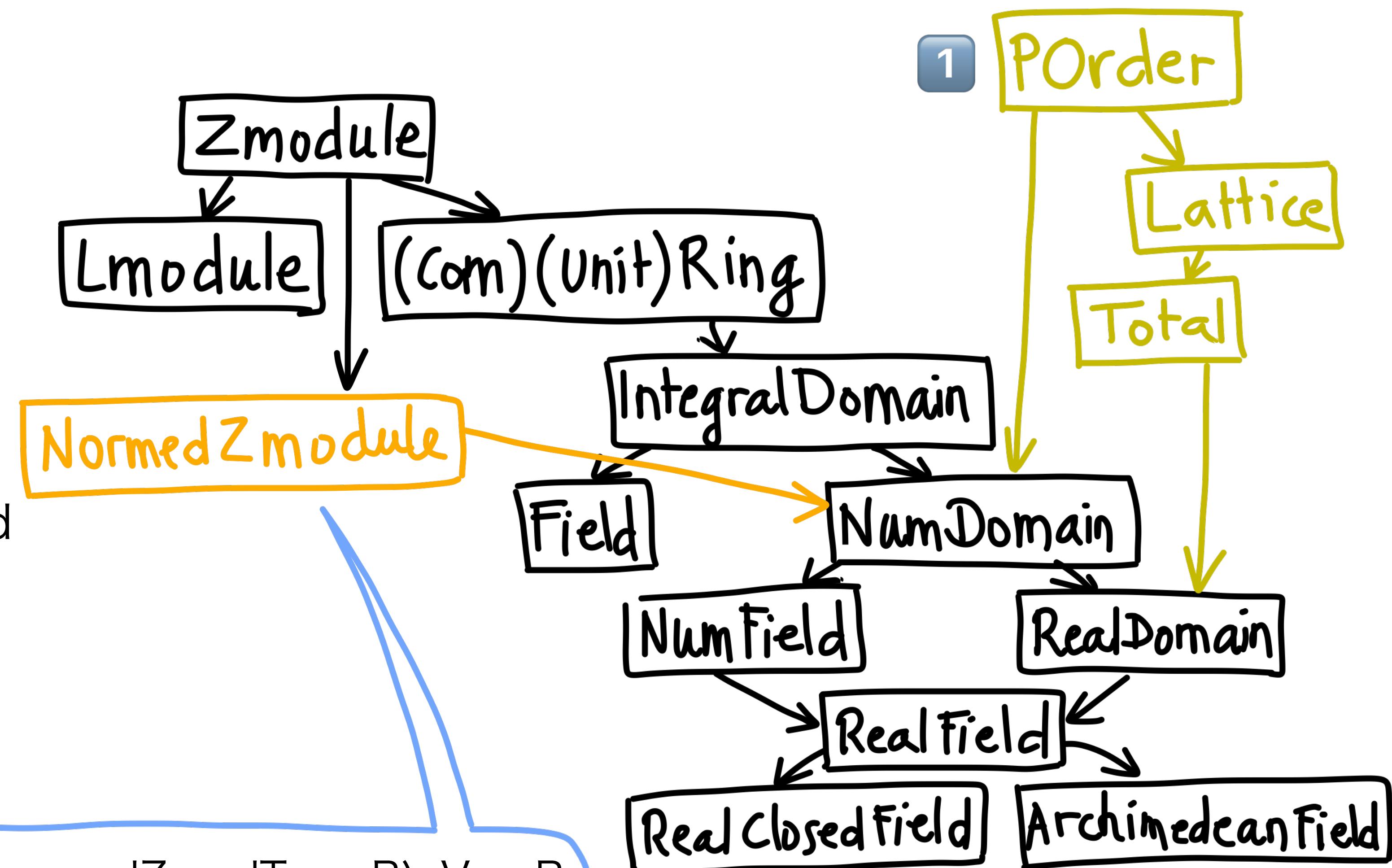
2 the class of NumDomain changed  
to include the mixin of  
NormedZmodule

3 Forgetful inheritance:

1. the class of NormedZmodule is  
parameterized by NumDomain and  
features the mixin of the norm

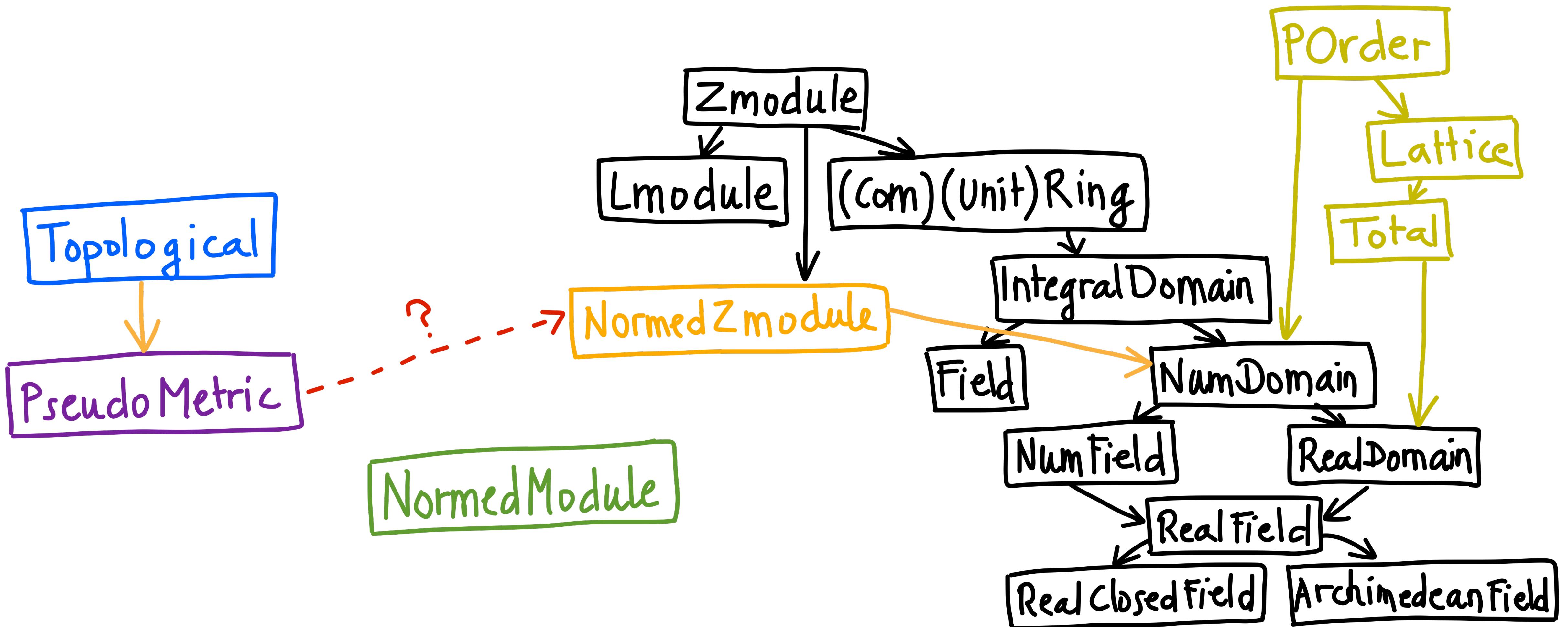
2. inheritance from NormedZmodule  
to NumDomain is declared

norm : forall (R : numDomainType) (V : normedZmodType R), V -> R



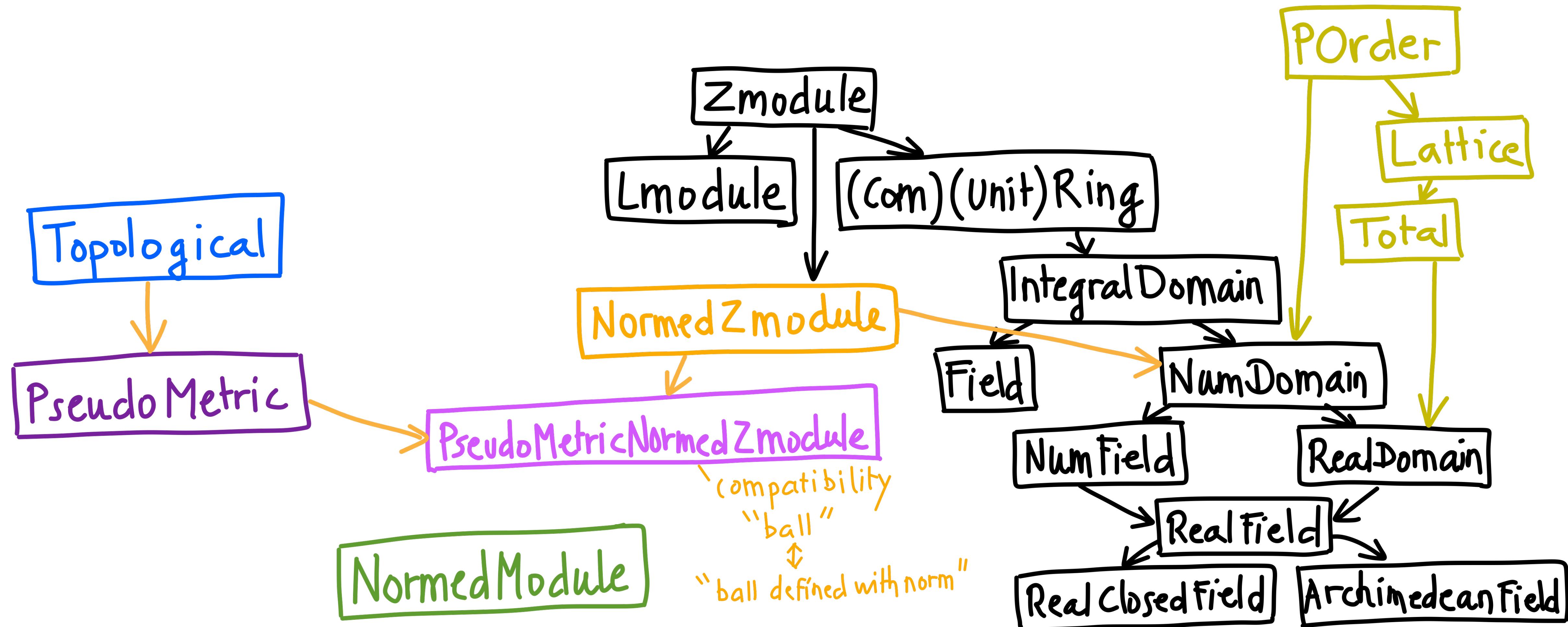
# How do we Integrate Normed Modules?

final step



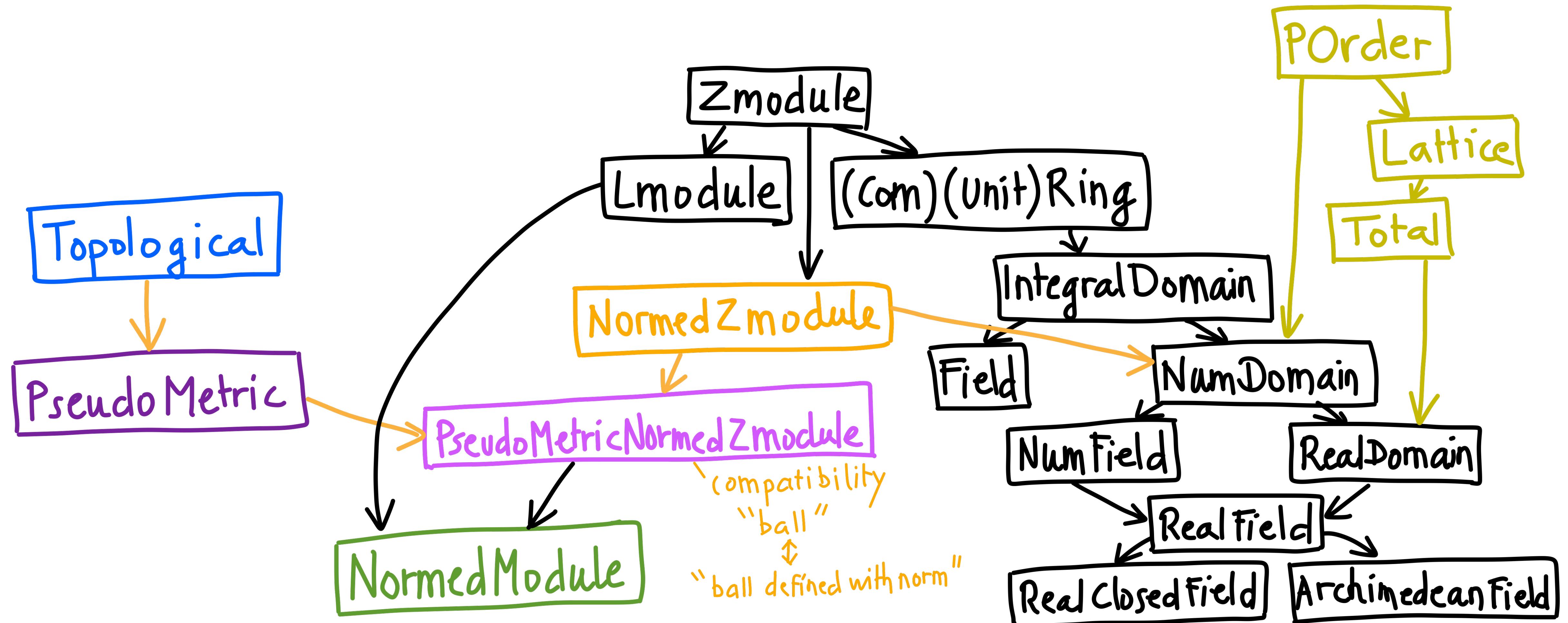
# Forgetful Inheritance

from normed modules to pseudometric spaces



# Forgetful Inheritance

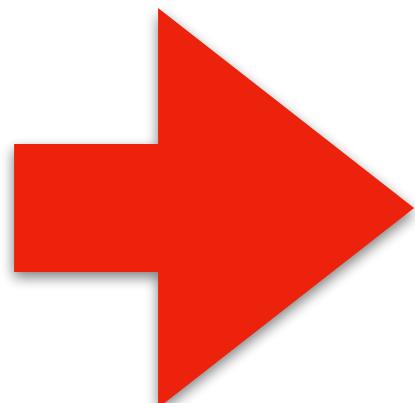
from normed modules to pseudometric spaces



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# Related Work

- HOL Light and Isabelle/HOL have more real and complex analysis
  - different foundations: we use dependent types here
- Forgetful inheritance was already mentioned
  - original description of packed classes [TPHOLs 2009] alludes to it
  - Buzzard et al. [CPP 2020] discuss a similar issue, on a specific example
- Packed classes + forgetful inheritance is verbose
  - Sakaguchi [IJCAR 2020] developed an automated checking tool ← on Friday
  - Cohen et al. [FSCD 2020] are working on automated generation ← tomorrow

# Conclusion

## summary of main contributions

1 forgetful inheritance using packed classes  
(several examples, comparison with type classes)

2 the MathComp-Analysis hierarchy  
(enhancement of MathComp,  
applications of forgetful inheritance)

Available: theories of Bachmann-Landau notation, of differentiability

